

**UNIVERSITY OF SWAZILAND
SUPPLEMENTARY EXAMINATION, JULY 2015**

Title of Paper : THEORY OF COMPUTATION

Course number: CS211

Time allowed : Three (3) hours.

Instructions : (1) Read all the questions in Section-A and Section-B from page one to page four before you start answering any question.

(2) Answer all questions in Section-A, and **any two** questions in section-B. Maximum mark is 100.

(3) Use correct notations and show all your work on the script.

This paper should not be opened until the invigilator has granted permission.

SECTION-A (Maximum marks 50)

Q1(a) (marks 6 + 6 +12). The following languages are given on symbol set $\{0, 1\}$. Assume that $u, v, w \in \{0, 1\}^*$.

(i). $L1 = \{uvw, |u| = 2 \text{ and } |v| = 1\}$

(ii). $L2 = \{0w0\} \cup \{1w1\}$

(iii). $L3 = \{w, (|w| \bmod 4 = 0)\}$

The following set of words is given -

$\{\lambda, 0, 1, 01, 001, 0100, 00011, 1111101, 0001111, 00000011, 001111011, 010101\}$

(a). From the above set, write all the words belonging to $L1$, all the words belonging to $L2$ and all the words belonging to $L3$.

(b). Write regular expressions representing $L1$, $L2$ and $L3$.

(c). Design three deterministic finite acceptors (**dfa's**) accepting $L1$, $L2$, and $L3$ respectively.

Q2 (marks 6 + 6 + 14). The following non deterministic finite acceptor (**nfa**) is given :

$M = (\{q0, q1, q2\}, \{0, 1\}, q0, \delta, \{q1, q2\})$, where the transitions are given as :

$$\delta(q0, 0) = \{q0, q1\} ;$$

$$\delta(q1, 1) = \{q0, q2\} ;$$

$$\delta(q2, 1) = \{q1, q2\} .$$

(a). Draw the transition digraph of the above **M**.

(b). Trace computations of all the words of L , where $L = \{111, 000 \text{ and } 010\}$.

(c). Find an equivalent **dfa** of **M** and write the state transition table of your **dfa**.

SECTION-B (Maximum marks 50)

Note: Answer **any two** questions in this section.

Q3(a) (marks 4+4+2). A context free grammar , $G = (\{S\}, \{a,b\}, S, P)$ where the set of productions P is given as

$$\{ S \rightarrow aS \mid aSbS \mid \lambda \}.$$

Find is the complexity of G . Write left most derivations and parse trees for $w_1 = aaaab$ and $w_2 = abab$. Taking example of w_1 , show that G is ambiguous.

Q3(b) (marks 9+6). Assuming n, m and $k \geq 0$. Find Context Free Grammars (CFG) G_1 and G_2 that generate the following languages. -

(i). $L(G_1) = \{a^{2n} b^i c^{2m}, \text{ such that } i = m + n\}$

(ii). $L(G_2) = \{a^n b^m c^k, \text{ such that } m = k\}.$

Write left most derivations for $w_1 = aab$, $w_2 = bcc$, $w_3 = aabbcc$ using G_1 and $w_4 = bbcc$, $w_5 = aabbcc$ and $w_6 = aabc$ using G_2 .

Include production number at each step of your derivation.

Q4(a) (marks 15). Design a deterministic pushdown automaton (**dpda**) to recognize the language –

$$L = \{w \in a^n b^m, \text{ such that } (n > 0) \text{ and } (m > n)\}$$

Clearly describe the functional steps as a numbered list of actions of your **dpda** of L . Write instantaneous descriptions for $w_1 = abb$ and $w_2 = aaabbbb$.

Q4(b) (marks 3 + 7). Design a non deterministic pushdown automaton (**npda**) to recognize the language generated by the grammar G , where –

$$G = (\{S, Z, B\}, \{a,b\}, S, P).$$

The set of productions P is

$$\begin{aligned} \{ S &\longrightarrow aZ \mid aB \mid a \\ B &\longrightarrow ZB \mid bB \mid b \\ Z &\longrightarrow aBZ \mid aZ \mid aB \mid a \} \end{aligned}$$

Q5 (marks 15 + 5 + 5). Write the design of a Turing Machine (TM) to compute –

$$F(x) = x \text{ div } 2.$$

Assume x to be a non zero positive integer in unary representation. Clearly write as to how your TM accomplishes the computations. Also write the instantaneous descriptions using the value of x as 11 and 11111 for your Turing Machine.

(End of Examination Paper)